INTERNATIONAL STANDARD

ISO/IEC 9798-3

Second edition 1998-10-15

AMENDMENT 1 2010-06-01

Information technology — Security techniques — Entity authentication —

Part 3:

Mechanisms using digital signature techniques

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Technologies de l'information — Techniques de sécurité — Is Authentification d'entité 2010

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Foreword

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Amendment 1 to ISO/IEC 9798-3:1998 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 27, 17 Security techniques.

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Information technology — Security techniques — Entity authentication —

Part 3: Mechanisms using digital signature techniques

AMENDMENT 1

Page 1, Clause 3

Replace the first paragraph of Clause 3 with the following:

For the purposes of this part of ISO/IEC 9798, the definitions and notation described in ISO/IEC 9798-1 and the following apply:

 I_A The identity of entity A, which is either A or CertA. **DREVEW**

 I_B The identity of entity B, which is either B or CertB en ai)

ResX The result of verifying entity X's public key or public key certificate.

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Page 5, 5.2.3

Add the following after 5.2.3:

6 Mechanisms involving an on-line trusted third party

6.1 Introduction

The authentication mechanisms in this clause require the two entities A and B to validate each other's public keys using an on-line trusted third party (with distinguishing identifier TP). This trusted third party shall possess reliable copies of the public keys of A and B. The entities A and B shall possess a reliable copy of the public key of TP.

This clause specifies two five pass authentication mechanisms, both of which achieve mutual authentication between entities *A* and *B*.

In the specification of the two mechanisms, the form of tokens and text fields follow the description given at the beginning of Clause 5, i.e. all paragraphs in Clause 5 before 5.1.

Implementations of the mechanisms shall use one of the signature schemes specified in ISO/IEC 14888 or ISO/IEC 9796.

6.2 Five pass authentication (initiated by A)

In this authentication mechanism, uniqueness/timeliness is controlled by generating and checking a random number (see Annex B of ISO/IEC 9798-1:1997).

This authentication mechanism is illustrated in Figure 6.



Figure 6 — Five pass authentication (initiated by A)

ISO/IEC 9798-3:1998/Amd 1:2010

The tokens shall be created according to one of the following two options: -baa5-4c1d-8966cfb6945ccab8/iso-iec-9798-3-1998-amd-1-2010

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 $\mathsf{Token}AB = \mathsf{Text9}||\mathsf{Res}A||sS_{7}(R_{B}||\mathsf{Res}A||\mathsf{Text5})||sS_{A}(R_{B}||R_{A}||B||A||\mathsf{Text8})$

 $\mathsf{Token}BA = R_A ||R_B||\mathsf{Text3}||sS_B(B||R_A||R_B||A||\mathsf{Text2})$

 $\mathsf{Token}TA = \mathsf{Res}A||\mathsf{Res}B||sS_{T}(R_{A}'||\mathsf{Res}B||\mathsf{Text6})||sS_{T}(R_{B}||\mathsf{Res}A||\mathsf{Text5})$

Option 2:

 $\mathsf{Token} AB = R'_A ||\mathsf{Text9}||\mathsf{Token} TA||sS_A(R_B||R_A||B||A||\mathsf{Text8})$

 $TokenBA = R_A ||R_B||Text3||sS_B(B||R_A||R_B||A||Text2)$

 $\mathsf{Token}TA = \mathsf{Res}A||\mathsf{Res}B||sS_T(R'_A||R_B||\mathsf{Res}A||\mathsf{Res}B||\mathsf{Text5})$

The values of the fields I_A , I_B , ResA, ResB, Status and Failure shall have the following forms:

 $I_A = A$ or CertA

 $I_B = B \text{ or Cert}B$

 $\text{Res}A = (\text{Cert}A||\text{Status}), (A||P_A) \text{ or Failure}$

 $\text{Res}B = (\text{Cert}B||\text{Status}), (B||P_B) \text{ or Failure}$

Status = True or False. The value of the field shall be set to False if the certificate is known to have been revoked; otherwise it shall be set to True.

Failure: ResX (where $X = \{A, B\}$) will be set to Failure if neither a public key nor a certificate of entity X can be found by TP.

In the mechanism, if *TP* knows the mapping between identity *X* and *P_X* (where *X* = {*A*, *B*}), then it shall set $I_X = X$; otherwise, it shall set $I_X = \text{Cert}X$, and *X* shall be set equal to the collection of distinguished identity fields in Cert*X*. If either *X* or Cert*X* is permitted to be used as an identity, then there should be a pre-arranged means to allow *TP* to distinguish the two types of identity indications. The value of Res*X* (where *X* = {*A*, *B*}) shall be determined according to Table 1.

Field	Choice 1	Choice 2
I _X	X	CertX
ResX	$(X P_X)$ or Failure	(CertX Status) or Failure

Table 1 — Value of ResX

The mechanism is performed as follows:

- 1) A sends a random number R_A , its identity I_A and, optionally, a text field Text1 to B.
- 2) B sends the token token BA and B to AARD PREVIEW
- 3) A sends a random number R_A , together with R_B , t_A , η_B and, optionally, a text field Text4 to TP.
- 4) On receipt of the message in Step (3) from A, TP performs the following steps. If I_A = A and I_B = B, TP retrieves P_A and P_B, If I_A = CertA and I_B = CertB₁, TP checks the validity of CertA and CertB. The process of certificate verification by TP may require protection from denial-of-service attacks. The specification of mechanisms to be used to provide such protection is outside of the scope of this part of ISO/IEC 9798.
- 5) Then *TP* sends Token*TA* and, optionally, a text field Text7 to *A*. The fields Res*A* and Res*B* in Token*TA* shall be: the certificates of *A* and *B* and their status, the distinguishing identifiers of *A* and *B* and their public keys, or an indication of Failure.
- 6) On receipt of the message in Step (5) from *TP*, *A* performs the following steps:
 - (i) Verify Token*TA* by checking the signature of *TP* contained in the token, and by checking that the random number R'_A , sent to *TP* in Step (3), is the same as the random number R'_A contained in the signed data of Token*TA*.
 - (ii) Retrieve the public key of *B* from the message, verify Token*BA* received in Step (2) by checking the signature of *B* contained in the token and checking that the value of identifier field (*A*) in the signed data of Token*BA* is equal to *A*'s distinguishing identifier, and then check that the random number R_A , sent to *B* in Step (1), is the same as the random number R_A contained in Token*BA*.
- 7) A sends TokenAB to B.
- 8) On receipt of the message in Step (7) from *A*, *B* performs the following steps:
 - (i) Verify Token *TA* by checking the signature of *TP* contained in the token, and by checking that the random number R_B , sent to *A* in Step (2), is the same as the random number R_B contained in the signed data of Token *TA*.

(ii) Retrieve the public key of *A* from the message, verify Token*AB* by checking the signature of *A* contained in the token and checking that the value of identifier field (*B*) in the signed data of Token*AB* is equal to *B*'s distinguishing identifier, and then check that the random number R_B contained in the signed data of Token*AB* is equal to the random number *R* sent to *A* in Step (2).

6.3 Five pass authentication (initiated by B)

In this authentication mechanism, uniqueness/timeliness is controlled by generating and checking a random number (see Annex B of ISO/IEC 9798-1).

This authentication mechanism is illustrated in Figure 7.



Figure 7 — Five pass authentication (initiated by *B*)

The tokens shall be created according to one of the following two options.

Option 1:

 $\mathsf{Token} AB = \mathsf{Text7} ||R_A||\mathsf{Res} A||sS_T(R_B||\mathsf{Res} A||\mathsf{Text3})||sS_A(R_B||R_A||B||A||\mathsf{Text6})$

 $TokenBA = R_A ||R_B||Text9||sS_B(A||R_A||R_B||B||Text8)$

 $\mathsf{Token}TA = \mathsf{Res}A||\mathsf{Res}B||sS_{T}(R_{A}'||\mathsf{Res}B||\mathsf{Text4})||sS_{T}(R_{B}||\mathsf{Res}A||\mathsf{Text3})$

Option 2:

 $TokenAB = R'_A ||Text7||TokenTA||sS_A(R_B||R_A||B||A||Text6)$

 $\mathsf{Token}BA = R_A ||R_B||\mathsf{Text9}||sS_B(R_A||R_B||A||B||\mathsf{Text8})$

Token*TA* = Res*A*||Res*B*|| $sS_T(R'_A||R_B||ResA||ResB||Text3)$

The values of the fields I_A, I_B, ResA, ResB, Status and Failure shall have the following forms:

 $I_A = A \text{ or Cert}A$

 $I_B = B$ or CertB

ResA = (CertA||Status), (A|| P_A) or Failure

 $\text{Res}B = (\text{Cert}B||\text{Status}), (B||P_B) \text{ or Failure}$

Status = True or False. The value of the field shall be set to False if the certificate is known to have been revoked; otherwise it shall be set to True.

Failure: ResY (where $Y = \{A, B\}$) will be set to Failure if neither a public key nor a certificate of entity Y can be found by TP.

In the mechanism, if *TP* knows the mapping between identity Y and P_Y (where Y = {A, B}), then it shall set $I_Y = Y$; otherwise, it shall set $I_Y = \text{Cert}Y$, and Y shall be set equal to the collection of distinguished identity fields in CertY. If either Y or CertY is permitted to be used as an identity, then there should be a pre-arranged means to allow *TP* to distinguish the two types of identity indications. The value of ResY (where Y = {A, B}) shall be determined according to Table 2.

Field	Choice 1	Choice 2
I _Y	Y	CertY
ResY	$(Y P_Y)$ or Failure	(CertY Status) or Failure
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Table 2 — Value of Res Y

The mechanism is performed as for standards.iteh.ai)

1) B sends a random number R_{B} its identity l_{B} and optionally, a text field Text1 to A.

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- 2) A sends a random number R'_{A} together with R_{B} , J_{AS} , J_{B} and, optionally, a text field Text2 to TP.
- 3) On receipt of the message in Step (2) from *A*, *TP* performs the following steps. If $I_A = A$ and $I_B = B$, *TP* retrieves P_A and P_B ; If $I_A = CertA$ and $I_B = CertB$, *TP* checks the validity of CertA and CertB. The process of certificate verification by *TP* may require protection from denial-of-service attacks. The specification of mechanisms to be used to provide such protection is outside of the scope of this part of ISO/IEC 9798.
- 4) Then *TP* sends Token*TA* and, optionally, a text field Text5 to *A*. The fields Res*A* and Res*B* in Token*TA* shall be: the certificates of *A* and *B* and their status, the distinguishing identifiers of *A* and *B* and their public keys or an indication of Failure.
- 5) A sends the token TokenAB and I_A to B.
- 6) On receipt of the message in Step (5) from *A*, *B* performs the following steps:
 - (i) Verify the signature of *TP* in Token*AB* by checking the signature of *TP* contained in the token, and by checking that the random number R_{B} , sent to *A* in Step (1), is the same as the random number R_{B} contained in the signed data of *TP* of Token*AB*.
 - (ii) Retrieve the public key of *A* from the message, verify Token*AB* by checking the signature of *A* contained in the token and checking that the value of identifier field (*B*) in the signed data of Token*AB* is equal to *B*'s distinguishing identifier, and then check that the random number R_B , sent to *A* in Step (1), is the same as the random number R_B contained in the signed data of *A* of Token*AB*.
- 7) B sends TokenBA to A.