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AMENDMENT 1
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**Information technology — Security
techniques — Random bit generation**

**AMENDMENT 1: Deterministic random
bit generation**

*Technologies de l'information — Techniques de sécurité —
Génération de bits aléatoires*

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AMENDEMENT 1: Génération déterministe de bits aléatoires

[ISO/IEC 18031:2011/Amd 1:2017](https://standards.iteh.ai/catalog/standards/sist/27cc2718-a470-4614-b5a8-88fd1ba017cf/iso-iec-18031-2011-amd-1-2017)

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Amendment 1 to ISO/IEC 18031-1:2011 was prepared by Joint Technical Committee ISO/IEC JTC 1, Information technology, Subcommittee SC 27, IT Security techniques.

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Add a new Annex K.

Annex K (informative) Example cases for MQ_DRBG

K.1 General

Annex K and its supporting files provides example cases for 14 settings listed in ISO/IEC 18031:2011, Table C.5. The supporting files are available at the following URL:

<http://standards.iso.org/iso/18031/>

In each of the 14 settings described in Annex K, the bitstring P provides a randomly selected system of multivariate quadratic equations that complies with the selection rules of C.5.2.5. The rank distribution resulting from the verification of rank conditions is detailed for each setting.

P is given in the format described in C.5.2.4 which is recalled below. Each example case also includes a sequence of consecutive input-output pairs for the Evaluate_MQ(...) function.

K.1.1 Format for representing field elements

Each system coefficient is an element of the binary field $\text{GF}(2^{\text{field_size}})$ and is a univariate polynomial over $\text{GF}(2)$ modulo the irreducible polynomial given in Table C.6. A field element is handled as a bitstring of field_size bits composed of its $\text{GF}(2)$ coefficients ordered by decreasing degree. For example, the polynomial $x^3 + x + 1$ in $\text{GF}(2^4)$ is represented as the bitstring 1011.

K.1.2 Format for representing a single multivariate quadratic equation

The quadratic system used in MQ_DRBG operates on $n = \text{state_length} / \text{field_size}$ variables and contains $n + m$ equations where $m = \text{block_length} / \text{field_size}$. A quadratic equation is written as the concatenation of its coefficients in lexicographic order and by decreasing degree. Therefore the coefficient of the monomial x_1x_1 appears first, followed by that of x_1x_2 and so forth, up to the coefficient of x_1x_n . The coefficient of the monomial x_2x_2 appears next, followed by that of x_2x_3 and so forth, until the last quadratic coefficient $x_{n-1}x_n$ is reached. Then linear coefficients appear, starting with the coefficient of the monomial x_1 and ending with that of x_n . When $\text{field_size} = 1$, the linear coefficients are omitted since the underlying field is $\text{GF}(2)$ and $x_i x_i = x_i$. The string ends with the constant coefficient of the quadratic equation.

K.1.3 Format for representing a complete system of quadratic equations

The quadratic system encoded into the bitstring P contains its $n + m$ quadratic equations concatenated in sequential order, starting with the coefficients of the first equation and ending with those of the $(n+m)$ -th equation. P is formed by the resulting bit string of length system_length .

K.1.4 Format for representing inputs and outputs

The input x to **Evaluate_MQ**(P, x) is a vector of n field elements and is given as a bitstring formed by concatenating their bitstring representations, starting with x_1 and ending with x_n . Similarly, the output $y || z$ is a vector of $n + m$ field elements represented in the same format.

K.1.5 Summary of example cases

Table K.1 summarizes the 14 example cases.

Table K.1 — Summary of example cases

requested_strength	block_length			
	112	128	192	256
80	K.2 Binary field GF(2) $n = 112$ $m = 112$ $min_weight = 4$ $min_rank \geq 106$	K.4 Binary field GF(2 ⁴) $n = 32$ $m = 32$ $min_weight = 5$ $min_rank \geq 30$	K.7 Binary field GF(2 ⁶) $n = 32$ $m = 32$ $min_weight = 5$ $min_rank \geq 30$	K.11 Binary field GF(2 ⁸) $n = 32$ $m = 32$ $min_weight = 5$ $min_rank \geq 30$
112	K.3 Binary field GF(2) $n = 120$ $m = 112$ $min_weight = 4$ $min_rank \geq 114$	K.5 Binary field GF(2) $n = 128$ $m = 128$ $min_weight = 4$ $min_rank \geq 122$	K.8 Binary field GF(2 ⁴) $n = 48$ $m = 48$ $min_weight = 5$ $min_rank \geq 44$	K.12 Binary field GF(2 ⁴) $n = 64$ $m = 64$ $min_weight = 5$ $min_rank \geq 60$
128		K.6 Same as K.5	K.9 Binary field GF(2 ³) $n = 64$ $m = 64$ $min_weight = 5$ $min_rank \geq 60$	K.13 Same as K.12
192			K.10 Binary field GF(2) $n = 200$ $m = 192$ $min_weight = 4$ $min_rank \geq 192$	K.14 Binary field GF(2 ²) $n = 128$ $m = 128$ $min_weight = 5$ $min_rank \geq 124$
256				K.15 Binary field GF(2) $n = 272$ $m = 256$ $min_weight = 4$ $min_rank \geq 264$

K.2 Example case for requested_strength = 80 and block_length = 112

K.2.1 System of multivariate quadratic equations

The bitstring P containing the system coefficients is provided in digital form in the file “coefficients-BL-112-Sec-80-F2.bin” in accordance with the format described in K.1.3.

The file contains 177212 bytes and its SHA-1 checksum in hexadecimal form is

95d78546df132777af932886a887da96aa9afa46

The ranks are distributed as follows:

106: 4561

108: 2213145

110: 58156950

112: 43613144

Sum: 103987800

K.2.2 Inputs and outputs

The bitstrings x , y and z are provided in digital form in accordance with the format described in K.1.4. Their hexadecimal values are:

$x = 00000000000000000000000000000001$

$y = \text{bb8cf180cbc3a6002c19c770ed0d}$

$z = 7847b864cfadf70fb359203e06d8$

$x = \text{bb8cf180cbc3a6002c19c770ed0d}$

$y = \text{a1e0811b5b7733113ca8e22dd2b1}$

$z = 57d27f7b0fc67aec0d5e8115cd93$

$x = \text{a1e0811b5b7733113ca8e22dd2b1}$

$y = 634ae5294dbc4cc79ce11cfb1d7$

$z = \text{c42c5cc5b5b61396df3fcf7a4e2b}$

$x = 634ae5294dbc4cc79ce11cfb1d7$

$y = 36701faea23130a0407a44f5e420$

$z = \text{bf3ddd3cbb141fcd96cbba66ebb9}$

$x = 36701faea23130a0407a44f5e420$

$y = 74b5baa1095f61eb6b15d317d5ed$

z = 7f4ad5787a0c5451bddcf2aef533

x = 74b5baa1095f61eb6b15d317d5ed

y = 62804addbe9da290c38e9de0fe71

z = 5f1f209b62cce21f75d9d03607a9

x = 62804addbe9da290c38e9de0fe71

y = 7d0892da52eed7facc377af1918f

z = 69d5bef53c03fa33a0273cf44c21

x = 7d0892da52eed7facc377af1918f

y = 8ee43a16842345d4cd182852cdea

z = ed479a677e6c2a3cffbbada0e765

x = 8ee43a16842345d4cd182852cdea

y = 2eb8cc9185445b2bab3f4b504aaf

z = 9407f0fe9393fa335051ac2bf414

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x = 2eb8cc9185445b2bab3f4b504aaf

y = 8deb10cb70bc3818209a576fb5cb

z = 6106cb8aa8e9a7de949a506b2278

K.3 Example case for *requested_strength* = 112 and *block_length* = 112

K.3.1 System of multivariate quadratic equations

The bitstring *P* containing the system coefficients is provided in digital form in the file “coefficients-BL-112-Sec-112-F2.bin” in accordance with the format described in K.1.3.

The file contains 210569 bytes and its SHA-1 checksum in hexadecimal form is

ae1c4ea33afc96e3aa421f6456055a7c7ee33989

The ranks are distributed as follows:

114: 5239

116: 2551294

118: 66936700

120: 50200265

Sum: 119693498

K.3.2 Inputs and outputs

The bitstrings x , y and z are provided in digital form in accordance with the format described in K.1.4. Their hexadecimal values are:

$x = 00000000000000000000000000000001$

$y = 46609cda28057a917a08b60a1d969d$

$z = a06fe3e456a8c24315dfde6088bd$

$x = 46609cda28057a917a08b60a1d969d$

$y = 37d12de7b69f2170ba8717e96f0f43$

$z = 8fb9899c9e2d4ef33056aadf946d$

$x = 37d12de7b69f2170ba8717e96f0f43$

$y = 463860297cec60797650c4897563d4$

$z = 89745528548d7bd3a2c9e5afd3fc$

$x = 463860297cec60797650c4897563d4$

$y = 6a4c5b16c156738e9b07c4c2c2818e$

$z = 5f9f14194e601f48657164f34e34$

$x = 6a4c5b16c156738e9b07c4c2c2818e$

$y = 289c50a28bb48a685703eb425597dd$

$z = c9dae7a3c32a01648a32d91b8728$

$x = 289c50a28bb48a685703eb425597dd$

$y = 4d96224af4aeaac54d8472374f645d$

$z = cf7a6cc73793049241497ee26603$

$x = 4d96224af4aeaac54d8472374f645d$

$y = df5ac81223125d967056d5dcdba088$

$z = 3d9741ec702076fe8473b7181aa9$

$x = df5ac81223125d967056d5dcdba088$

$y = 41a1df8cc57c402f520d671464b728$

$z = 285d6b741e417e417b9f8fa87356$

$x = 41a1df8cc57c402f520d671464b728$

$y = 0af3539a48bc07e3afb00d3c529ff5$

$z = e6d4d36dcc2cca4826b94e76be10$

$x = 0af3539a48bc07e3afb00d3c529ff5$

$y = e2f7d8f01d2ae145a643b9351ada76$

$z = 29bdd54840cf84027f20e48ce195$

K.4 Example case for *requested_strength* = 80 and *block_length* = 128

K.4.1 System of multivariate quadratic equations

The bitstring P containing the system coefficients is provided in digital form in the file “coefficients-BL-128-Sec-80-F16.bin” in accordance with the format described in K.1.3.

The file contains 17952 bytes and its SHA-1 checksum in hexadecimal form is

d6614e19bd953ca88ff49f016b80f5ac17b7dab1

<https://standards.iteh.ai/catalog/standards/sist/27cc2718-a470-4614-b5a8-88fd1ba017cf/iso-iec-18031-2011-amd-1-2017>

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The ranks are distributed as follows:

30: 520948

32: 7782684

Sum: 8303632

K.4.2 Inputs and outputs

The bitstrings x , y and z are provided in digital form in accordance with the format described in K.1.4. Their hexadecimal values are:

$x = 00000000000000000000000000000001$

$y = f719e81ed992ca7c793258b5251d0534$

$z = 66092272f74a85ecaef639d78ed9831f$

$x = f719e81ed992ca7c793258b5251d0534$

$y = 37614b89b9bbd6eea4560ecb3bdb8807$

$z = 96b4c1aeb27aa47fbc7a3b1464343736$

x = 37614b89b9bbd6eea4560ecb3bdb8807
 y = 136bf7d8fbcabd37a2baa321a5d94f7
 z = 29141359d8099496eaf84ae3d863591a

x = 136bf7d8fbcabd37a2baa321a5d94f7
 y = bc6316205ac244b4fc8dcee70f423874
 z = d8005ccefafa012118820cf02c9eb4328d

x = bc6316205ac244b4fc8dcee70f423874
 y = 64d8adbf03a6418fa549f235e5f84bcd
 z = 9c0aad312ef00336d0f055e81f2b3677

x = 64d8adbf03a6418fa549f235e5f84bcd
 y = 3ac1c733b68ca734550343d950649d5a
 z = 1f07210c4a6d4fd784ee0f9f9789c5ab

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x = 3ac1c733b68ca734550343d950649d5a
 y = 1a22cbbe771e641373700306718dbf6e
 z = ba8064102a7e8d714e92e0dfddfbe607

x = 1a22cbbe771e641373700306718dbf6e
 y = fa2eabf2c9794f6b9bac6561409aab0d
 z = 7e2bae34daaf284557bbe5ae48e54d26

x = fa2eabf2c9794f6b9bac6561409aab0d
 y = 46f6f74d23504a64565b2c35cd0036df
 z = c6285e77cbf16150457d03bfc6015ef7

x = 46f6f74d23504a64565b2c35cd0036df
 y = 729bc30c32fd7fec1ccb95bc4aabfa27
 z = 963bda8ab7dc84ee2dd5a60a9c4392cd

K.5 Example case for *requested_strength* = 112 and *block_length* = 128